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NEW RESULTS ON A GENERALIZED COUPON COLLECTOR PROBLEM USING MARKOV CHAINS

EMMANUELLE ANCEAUME, YANN BUSNEL AND BRUNO SERICOLA

ABSTRACT. We study in this paper a generalized coupon collector problem, which consists in determining the distribution and the moments of the time needed to collect a given number of distinct coupons that are drawn from a set of coupons with an arbitrary probability distribution. We suppose that a special coupon called the null coupon can be drawn but never belongs to any collection. In this context, we obtain expressions of the distribution and the moments of this time. We also prove that the almost-uniform distribution, for which all the non-null coupons have the same drawing probability, is the distribution which minimizes the expected time to get a fixed subset of distinct coupons. This optimization result is extended to the complementary distribution of that time when the full collection is considered, proving by the way this well-known conjecture. Finally, we propose a new conjecture which expresses the fact that the almost-uniform distribution should minimize the complementary distribution of the time needed to get any fixed number of distinct coupons.

Keywords: Coupon collector problem; Minimization; Markov chains

1. INTRODUCTION

The coupon collector problem is an old problem which consists in evaluating the time needed to get a collection of different objects drawn randomly using a given probability distribution. This problem has given rise to a lot of attention from researchers in various fields since it has applications in many scientific domains including computer science and optimization.

More formally, consider a set of n coupons which are drawn randomly one by one, with replacement, coupon i being drawn with probability p_i . The classical coupon collector problem is to determine the expectation or the distribution of the number of coupons that need to be drawn from the set of n coupons to obtain the full collection of the n coupons. A large number of papers have been devoted to the analysis of asymptotics and limit distributions of this distribution when n tends to infinity, see [3] or [6] and the references therein. In [2], the authors obtain some new formulas concerning this distribution and they also provide simulation techniques to compute it as well as analytic bounds of it.

We consider in this paper several generalizations of this problem. A first generalization is the analysis, for $c \leq n$, of the number $T_{c,n}$ of coupons that need to be drawn, with replacement, to collect c different coupons from set $\{1, 2, \dots, n\}$. With this notation, the number of coupons that need to be drawn from this set to obtain the full collection is $T_{n,n}$. If a coupon is drawn at each discrete time $1, 2, \dots$ then $T_{c,n}$ is the time needed to obtain c different coupons also called the waiting time to obtain c different coupons. This problem has been considered in [7] in the case where the drawing probability distribution is uniform.

In a second generalization, we assume that $p = (p_1, \dots, p_n)$ is not necessarily a probability distribution, *i.e.*, we suppose that $p_1 + \dots + p_n \leq 1$ and we define $p_0 = 1 - (p_1 + \dots + p_n)$. This means that there is a null coupon, denoted by 0, which is drawn with probability p_0 , but which does not belong to the collection. In this context, the problem is to determine the distribution of the number $T_{c,n}$ of coupons that need to be drawn from set $\{0, 1, 2, \dots, n\}$, with replacement, till one first obtains a collection composed of c different coupons, $1 \leq c \leq n$, among $\{1, \dots, n\}$. These generalizations are motivated by the analysis of streaming algorithms in network monitoring applications as presented in Section 7.

The distribution of $T_{c,n}$ is obtained using Markov chains in Section 2, in which we moreover show that this distribution leads to new combinatorial identities. This result is used to get an expression of $T_{c,n}(v)$ when the drawing distribution is the almost-uniform distribution denoted by v and defined by $v = (v_1, \dots, v_n)$ with $v_i = (1 - v_0)/n$, where $v_0 = 1 - (v_1 + \dots + v_n)$. Expressions of the moments of $T_{c,n}(p)$ are given in Section 3, where we show that the limit of $E(T_{c,n}(p))$ is equal to c when n tends to infinity. We show in

Section 4 that the almost-uniform distribution v and the uniform distribution u minimize the expected value $E(T_{c,n}(p))$. We prove in Section 5 that the tail distribution of $T_{c,n}$ is minimized over all the p_1, \dots, p_n by the almost-uniform distribution and by the uniform distribution. This result was expressed as a conjecture in the case where $p_0 = 0$, *i.e.*, when $p_1 + \dots + p_n = 1$, in several papers like [1] for instance, from which the idea of the proof comes from. We propose in Section 6 a new conjecture which consists in showing that the distributions v and u minimize the tail distribution of $T_{c,n}(p)$. This conjecture is motivated by the fact that it is true for $c = 1$ and $c = n$ as shown in Section 5, and we show that it is also true for $c = 2$. It is moreover true for the expected value $E(T_{c,n}(p))$ as shown in Section 4.

2. DISTRIBUTION OF $T_{c,n}$

Recall that $T_{c,n}$ is the number of coupons that need to be drawn from set $\{0, 1, 2, \dots, n\}$, with replacement, till one first obtains a collection with c different coupons, $1 \leq c \leq n$, among $\{1, \dots, n\}$, where coupon i is drawn with probability p_i , $i = 0, 1, \dots, n$.

To obtain the distribution of $T_{c,n}$, we consider the discrete-time Markov chain $X = \{X_m, m \geq 0\}$ that represents the collection obtained after having drawn m coupons. The state space of X is

$$S_n = \{J \subseteq \{1, \dots, n\}\}$$

with transition probability matrix, denoted by Q , given, for every $J, H \in S_n$, by

$$Q_{J,H} = \begin{cases} p_\ell & \text{if } H \setminus J = \{\ell\} \\ p_0 + P_J & \text{if } J = H \\ 0 & \text{otherwise,} \end{cases}$$

where, for every $J \in S_n$, P_J is given by

$$(1) \quad P_J = \sum_{j \in J} p_j,$$

with $P_\emptyset = 0$. It is easily checked that Markov chain X is acyclic, *i.e.*, it has no cycle of length greater than 1, and that all the states are transient, except state $\{1, \dots, n\}$ which is absorbing. We introduce the partition $(S_{0,n}, S_{1,n}, \dots, S_{n,n})$ of S_n , where $S_{i,n}$ is defined, for $i = 0, \dots, n$, by

$$(2) \quad S_{i,n} = \{J \subseteq \{1, \dots, n\} \mid |J| = i\}.$$

Note that we have $S_{0,n} = \{\emptyset\}$,

$$|S_n| = 2^n \text{ and } |S_{i,n}| = \binom{n}{i}.$$

Assuming that $X_0 = \emptyset$ with probability 1, the random variable $T_{c,n}$ can then be defined, for every $c = 1, \dots, n$, by

$$T_{c,n} = \inf\{m \geq 0 \mid X_m \in S_{c,n}\}.$$

The distribution of $T_{c,n}$ is obtained in Theorem 2 using the Markov property and the following lemma.

For every $n \geq 1$, $\ell = 1, \dots, n$ and $i = 0, \dots, n$, we define the set $S_{i,n}(\ell)$ by

$$S_{i,n}(\ell) = \{J \subseteq \{1, \dots, n\} \setminus \{\ell\} \mid |J| = i\}.$$

Lemma 1. *For every $n \geq 1$, for every $k \geq 0$, for all positive real numbers y_1, \dots, y_n , for every $i = 1, \dots, n$ and all real number $a \geq 0$, we have*

$$\sum_{\ell=1}^n y_\ell \sum_{J \in S_{i-1,n}(\ell)} (a + y_\ell + Y_J)^k = \sum_{J \in S_{i,n}} Y_J (a + Y_J)^k,$$

where $Y_J = \sum_{j \in J} y_j$ and $Y_\emptyset = 0$.

Proof. For $n = 1$, since $S_{0,1}(1) = \emptyset$, the left hand side is equal to $y_1(a + y_1)^k$ and since $S_{1,1} = \{1\}$, the right hand side is also equal to $y_1(a + y_1)^k$. Suppose that the result is true for integer $n - 1$, with $n \geq 2$, i.e., suppose that for every $k \geq 0$, for all positive real numbers y_1, \dots, y_{n-1} , for every $i = 1, \dots, n - 1$ and for all real number $a \geq 0$, we have

$$\sum_{\ell=1}^{n-1} y_\ell \sum_{J \in S_{i-1,n-1}(\ell)} (a + y_\ell + Y_J)^k = \sum_{J \in S_{i,n-1}} Y_J(a + Y_J)^k.$$

We then have

$$\sum_{\ell=1}^n y_\ell \sum_{J \in S_{i-1,n}(\ell)} (a + y_\ell + Y_J)^k = \sum_{\ell=1}^{n-1} y_\ell \sum_{J \in S_{i-1,n}(\ell)} (a + y_\ell + Y_J)^k + y_n \sum_{J \in S_{i-1,n}(n)} (a + y_n + Y_J)^k.$$

Since $S_{i-1,n}(n) = S_{i-1,n-1}$, we get

$$\sum_{\ell=1}^n y_\ell \sum_{J \in S_{i-1,n}(\ell)} (a + y_\ell + Y_J)^k = \sum_{\ell=1}^{n-1} y_\ell \sum_{J \in S_{i-1,n}(\ell)} (a + y_\ell + Y_J)^k + y_n \sum_{J \in S_{i-1,n-1}} (a + y_n + Y_J)^k.$$

For $\ell = 1, \dots, n - 1$, the set $S_{i-1,n}(\ell)$ can be partitioned into two subsets $S'_{i-1,n}(\ell)$ and $S''_{i-1,n}(\ell)$ defined by

$$S'_{i-1,n}(\ell) = \{J \subseteq \{1, \dots, n\} \setminus \{\ell\} \mid |J| = i - 1 \text{ and } n \in J\}$$

and

$$S''_{i-1,n}(\ell) = \{J \subseteq \{1, \dots, n\} \setminus \{\ell\} \mid |J| = i - 1 \text{ and } n \notin J\}.$$

Since $S''_{i-1,n}(\ell) = S_{i-1,n-1}(\ell)$, the previous relation becomes

$$\begin{aligned} & \sum_{\ell=1}^n y_\ell \sum_{J \in S_{i-1,n}(\ell)} (a + y_\ell + Y_J)^k \\ &= \sum_{\ell=1}^{n-1} y_\ell \left[\sum_{J \in S_{i-1,n-1}(\ell)} (a + y_\ell + Y_J)^k + \sum_{J \in S'_{i-1,n}(\ell)} (a + y_\ell + Y_J)^k \right] + y_n \sum_{J \in S_{i-1,n-1}} (a + y_n + Y_J)^k \\ &= \sum_{\ell=1}^{n-1} y_\ell \sum_{J \in S_{i-1,n-1}(\ell)} (a + y_\ell + Y_J)^k + \sum_{\ell=1}^{n-1} y_\ell \sum_{J \in S_{i-2,n-1}(\ell)} (a + y_n + y_\ell + Y_J)^k + y_n \sum_{J \in S_{i-1,n-1}} (a + y_n + Y_J)^k. \end{aligned}$$

The recurrence hypothesis can be applied for both the first and the second terms. For the second term, the constant a is replaced by the constant $a + y_n$. We thus obtain

$$\begin{aligned} & \sum_{\ell=1}^n y_\ell \sum_{J \in S_{i-1,n}(\ell)} (a + y_\ell + Y_J)^k \\ &= \sum_{J \in S_{i,n-1}} Y_J(a + Y_J)^k + \sum_{J \in S_{i-1,n-1}} Y_J(a + y_n + Y_J)^k + y_n \sum_{J \in S_{i-1,n-1}} (a + y_n + Y_J)^k \\ &= \sum_{J \in S_{i,n-1}} Y_J(a + Y_J)^k + \sum_{J \in S_{i-1,n-1}} (y_n + Y_J)(a + y_n + Y_J)^k \\ &= \sum_{J \in S_{i,n-1}} Y_J(a + Y_J)^k + \sum_{J \in S'_{i,n}} Y_J(a + Y_J)^k, \end{aligned}$$

where $S'_{i,n} = \{J \subseteq \{1, \dots, n\} \mid |J| = i \text{ and } n \in J\}$.

Consider the set $S''_{i,n} = \{J \subseteq \{1, \dots, n\} \mid |J| = i \text{ and } n \notin J\}$. The sets $S'_{i,n}$ and $S''_{i,n}$ form a partition of $S_{i,n}$ and since $S''_{i,n} = S_{i,n-1}$, we get

$$\begin{aligned} \sum_{\ell=1}^n y_{\ell} \sum_{J \in S_{i-1,n}(\ell)} (a + y_{\ell} + Y_J)^k &= \sum_{J \in S_{i,n-1}} Y_J(a + Y_J)^k + \sum_{J \in S'_{i,n}} Y_J(a + Y_J)^k \\ &= \sum_{J \in S''_{i,n}} Y_J(a + Y_J)^k + \sum_{J \in S'_{i,n}} Y_J(a + Y_J)^k \\ &= \sum_{J \in S_{i,n}} Y_J(a + Y_J)^k, \end{aligned}$$

which completes the proof. \square

In the following we will use the fact that the distribution of $T_{c,n}$ depends on the vector $p = (p_1, \dots, p_n)$, so we will use the notation $T_{c,n}(p)$ instead of $T_{c,n}$, meaning by the way that vector p is of dimension n . We will also use the notation

$$p_0 = 1 - \sum_{i=1}^n p_i.$$

Finally, for $\ell = 1, \dots, n$, the notation $p^{(\ell)}$ will denote the vector p in which the entry p_{ℓ} has been removed, that is $p^{(\ell)} = (p_i)_{1 \leq i \leq n, i \neq \ell}$. The dimension of $p^{(\ell)}$, which is $n - 1$ here, is not specified but will be clear by the context of its use. We are now able to prove the following result.

Theorem 2. *For every $n \geq 1$ and $c = 1, \dots, n$, we have, for every $k \geq 0$,*

$$(3) \quad \mathbb{P}\{T_{c,n}(p) > k\} = \sum_{i=0}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \sum_{J \in S_{i,n}} (p_0 + P_J)^k,$$

where P_J is given by (1).

Proof. Relation (3) is true for $c = 1$ since in this case we have

$$\mathbb{P}\{T_{1,n}(p) > k\} = p_0^k.$$

So we suppose now that $n \geq 2$ and $c = 2, \dots, n$.

Since $X_0 = \emptyset$, conditioning on X_1 and using the Markov property, see for instance [8], we get for $k \geq 1$,

$$(4) \quad \mathbb{P}\{T_{c,n}(p) > k\} = p_0 \mathbb{P}\{T_{c,n}(p) > k-1\} + \sum_{\ell=1}^n p_{\ell} \mathbb{P}\{T_{c-1,n-1}(p^{(\ell)}) > k-1\}.$$

We now proceed by recurrence over index k . Relation (3) is true for $k = 0$ since it is well-known that

$$(5) \quad \sum_{i=0}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \binom{n}{i} = 1.$$

Relation (3) is also true for $k = 1$ since on the one hand $\mathbb{P}\{T_{c,n}(p) > 1\} = 1$ and on the other hand, using Relation (4), we have

$$\begin{aligned} \mathbb{P}\{T_{c,n}(p) > 1\} &= p_0 \mathbb{P}\{T_{c,n}(p) > 0\} + \sum_{\ell=1}^n p_{\ell} \mathbb{P}\{T_{c-1,n-1}(p^{(\ell)}) > 0\} \\ &= p_0 + \sum_{\ell=1}^n p_{\ell} \\ &= 1. \end{aligned}$$

Suppose now that Relation (3) is true for integer $k - 1$, that is, suppose that we have

$$P\{T_{c,n}(p) > k - 1\} = \sum_{i=0}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \sum_{J \in S_{i,n}} (p_0 + P_J)^{k-1}.$$

Using (4) and the recurrence relation, we have

$$\begin{aligned} P\{T_{c,n}(p) > k\} &= p_0 \sum_{i=0}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \sum_{J \in S_{i,n}} (p_0 + P_J)^{k-1} \\ &\quad + \sum_{\ell=1}^n p_\ell \sum_{i=0}^{c-2} (-1)^{c-2-i} \binom{n-i-2}{n-c} \sum_{J \in S_{i,n}(\ell)} (p_0 + p_\ell + P_J)^{k-1}. \end{aligned}$$

Using the change of variable $i := i - 1$ in the second sum, we obtain

$$\begin{aligned} P\{T_{c,n}(p) > k\} &= \sum_{i=0}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} p_0 \sum_{J \in S_{i,n}} (p_0 + P_J)^{k-1} \\ &\quad + \sum_{i=1}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \sum_{\ell=1}^n p_\ell \sum_{J \in S_{i-1,n}(\ell)} (p_0 + p_\ell + P_J)^{k-1} \\ &= \sum_{i=1}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \left[p_0 \sum_{J \in S_{i,n}} (p_0 + P_J)^{k-1} + \sum_{\ell=1}^n p_\ell \sum_{J \in S_{i-1,n}(\ell)} (p_0 + p_\ell + P_J)^{k-1} \right] \\ &\quad + (-1)^{c-1} \binom{n-1}{n-c} p_0^k. \end{aligned}$$

From Lemma 1, we have

$$\sum_{\ell=1}^n p_\ell \sum_{J \in S_{i-1,n}(\ell)} (p_0 + p_\ell + P_J)^{k-1} = \sum_{J \in S_{i,n}} P_J (p_0 + P_J)^{k-1},$$

that is

$$\begin{aligned} P\{T_{c,n}(p) > k\} &= (-1)^{c-1} \binom{n-1}{n-c} p_0^k + \sum_{i=1}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \sum_{J \in S_{i,n}} (p_0 + P_J)^k \\ &= \sum_{i=0}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \sum_{J \in S_{i,n}} (p_0 + P_J)^k, \end{aligned}$$

which completes the proof. \square

This theorem also shows, as expected, that the function $P\{T_{c,n}(p) > k\}$, as a function of p , is symmetric, which means that it has the same value for any permutation of the entries of p . As a corollary, we obtain the following combinatorial identities.

Corollary 3. *For every $c \geq 1$, for every $n \geq c$ and for all $p_1, \dots, p_n \in (0, 1)$ such that $p_1 + \dots + p_n = 1$, we have*

$$\sum_{i=0}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \sum_{J \in S_{i,n}} (p_0 + P_J)^{k-1} = 1, \text{ for } k = 0, 1, \dots, c-1.$$

Proof. The random variable $T_{c,n}$ takes its values on the set $\{c, c+1, \dots\}$, so we have

$$P\{T_{c,n} > k\} = 1, \text{ for } k = 0, 1, \dots, c-1,$$

which completes the proof thanks to Theorem 2. \square

For every $n \geq 1$ and for every $v_0 \in [0, 1]$, we define the vector $v = (v_1, \dots, v_n)$ by $v_i = (1 - v_0)/n$. We will refer it to as the almost-uniform distribution. We then have, from (3),

$$\mathbb{P}\{T_{c,n}(v) > k\} = \sum_{i=0}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \binom{n}{i} \left(v_0 \left(1 - \frac{i}{n} \right) + \frac{i}{n} \right)^k.$$

We denote by $u = (u_1, \dots, u_n)$ the uniform distribution defined by $u_i = 1/n$. It is equal to v when $v_0 = 0$. The dimensions of u and v are specified by the context.

3. MOMENTS OF $T_{c,n}$

For $r \geq 1$, the r th moment of $T_{c,n}(p)$ is defined by

$$\mathbb{E}(T_{c,n}^r(p)) = \sum_{k=1}^{\infty} k^r \mathbb{P}\{T_{c,n}(p) = k\}.$$

It can be obtained in function of the tail distribution of $T_{c,n}(p)$ by writing

$$\begin{aligned} \mathbb{E}(T_{c,n}^r(p)) &= \sum_{k=1}^{\infty} k^r \mathbb{P}\{T_{c,n}(p) = k\} \\ &= \sum_{k=1}^{\infty} k^r \mathbb{P}\{T_{c,n}(p) > k-1\} - \sum_{k=1}^{\infty} k^r \mathbb{P}\{T_{c,n}(p) > k\} \\ &= \sum_{k=0}^{\infty} ((k+1)^r - k^r) \mathbb{P}\{T_{c,n}(p) > k\} \\ &= \sum_{\ell=0}^{r-1} \binom{r}{\ell} \sum_{k=0}^{\infty} k^{\ell} \mathbb{P}\{T_{c,n}(p) > k\}. \end{aligned}$$

We easily get the first and second moments of $T_{c,n}(p)$, by taking $r = 1$ and $r = 2$ respectively, that is

$$(6) \quad \mathbb{E}(T_{c,n}(p)) = \sum_{k=0}^{\infty} \mathbb{P}\{T_{c,n}(p) > k\} = \sum_{i=0}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \sum_{J \in S_{i,n}} \frac{1}{1 - (p_0 + P_J)}$$

and

$$\mathbb{E}(T_{c,n}^2(p)) = \mathbb{E}(T_{c,n}(p)) + 2 \sum_{k=1}^{\infty} k \mathbb{P}\{T_{c,n}(p) > k\} = \sum_{i=0}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \sum_{J \in S_{i,n}} \frac{1 + 2(p_0 + P_J)}{[1 - (p_0 + P_J)]^2}.$$

The expected value given by (6) has been obtained in [4] in the particular case where $p_0 = 0$.

When the drawing probabilities are given by the almost-uniform distribution v , we get

$$\begin{aligned} \mathbb{E}(T_{c,n}(v)) &= \frac{1}{1-v_0} \sum_{i=0}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \binom{n}{i} \frac{n}{n-i} \\ &= \frac{1}{1-v_0} \mathbb{E}(T_{c,n}(u)). \end{aligned}$$

Using the following two relations

$$\binom{n}{i} = \binom{n-1}{i} + \binom{n-1}{i-1} 1_{\{i \geq 1\}} \quad \text{and} \quad \binom{n-1}{i} \frac{n}{n-i} = \binom{n}{i},$$

where 1_A is the indicator function of set A , we get

$$\begin{aligned} \mathbb{E}(T_{c,n}(u)) &= \sum_{i=0}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \binom{n}{i} \frac{n}{n-i} \\ &= \sum_{i=0}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \binom{n}{i} + \sum_{i=1}^{c-1} (-1)^{c-1-i} \binom{n-i-1}{n-c} \binom{n-1}{i-1} \frac{n}{n-i}. \end{aligned}$$

From Relation (5), the first sum is equal to 1. Using the change of variable $i := i + 1$ in the second sum, we obtain

$$\begin{aligned} \mathbb{E}(T_{c,n}(u)) &= 1 + \sum_{i=0}^{c-2} (-1)^{c-2-i} \binom{n-i-2}{n-c} \binom{n-1}{i} \frac{n}{n-i+1} \\ (7) \quad &= 1 + \frac{n}{n-1} \mathbb{E}(T_{c-1,n-1}(u)). \end{aligned}$$

Note that the dimension of the uniform distribution in the left hand side is equal to n and the one in the right hand side is equal to $n - 1$. Since $\mathbb{E}(T_{1,n}(u)) = 1$, we obtain

$$(8) \quad \mathbb{E}(T_{c,n}(u)) = n(H_n - H_{n-c}) \text{ and } \mathbb{E}(T_{c,n}(v)) = \frac{n(H_n - H_{n-c})}{1 - v_0},$$

where H_ℓ is the ℓ th harmonic number defined by $H_0 = 0$ and, for $\ell \geq 1$,

$$H_\ell = \sum_{i=1}^{\ell} 1/i.$$

We deduce easily from (7) that, for every $c \geq 1$, we have

$$\lim_{n \rightarrow \infty} \mathbb{E}(T_{c,n}(u)) = c \text{ and } \lim_{n \rightarrow \infty} \mathbb{E}(T_{c,n}(v)) = \frac{c}{1 - v_0}.$$

In the next section we show that, when p_0 is fixed, the minimum value of $\mathbb{E}(T_{c,n}(p))$ is reached when $p = v$, with $v_0 = p_0$.

4. DISTRIBUTION MINIMIZING $\mathbb{E}(T_{c,n}(p))$

The following lemma will be used to prove the next theorem.

Lemma 4. *For every $n \geq 1$ and $r_1, \dots, r_n \in (0, 1)$ with $r_1 + \dots + r_n = 1$, we have*

$$\sum_{\ell=1}^n \frac{1}{r_\ell} \geq n^2.$$

Proof. We proceed by recurrence. The result is clearly true for $n = 1$. Suppose that the result is true for integer $n - 1$, $n \geq 2$. We then have

$$\sum_{\ell=1}^n \frac{1}{r_\ell} = \frac{1}{r_n} + \sum_{\ell=1}^{n-1} \frac{1}{r_\ell} = \frac{1}{r_n} + \frac{1}{1 - r_n} \sum_{\ell=1}^{n-1} \frac{1}{h_\ell},$$

where h_ℓ is given, for $\ell = 1, \dots, n - 1$, by

$$h_\ell = \frac{r_\ell}{1 - r_n}.$$

Since $h_1 + \dots + h_{n-1} = 1$, we get, using the recurrence hypothesis,

$$\sum_{\ell=1}^n \frac{1}{r_\ell} \geq \frac{1}{r_n} + \frac{(n-1)^2}{1 - r_n} = \frac{(nr_n - 1)^2}{r_n(1 - r_n)} + n^2 \geq n^2,$$

which completes the proof. ■

Theorem 5. *For every $n \geq 1$, for every $c = 1, \dots, n$ and $p = (p_1, \dots, p_n) \in (0, 1)^n$ with $p_1 + \dots + p_n \leq 1$, we have*

$$\mathbb{E}(T_{c,n}(p)) \geq \mathbb{E}(T_{c,n}(v)) \geq \mathbb{E}(T_{c,n}(u)),$$

where $v = (v_1, \dots, v_n)$ with $v_i = (1 - p_0)/n$ and $p_0 = 1 - (p_1 + \dots + p_n)$ and where $u = (1/n, \dots, 1/n)$.

Proof. The second inequality comes from (8).

Defining $v_0 = 1 - (v_1 + \dots + v_n)$, we have $v_0 = p_0$. For $c = 1$, the result is trivial since we have from Relation (6)

$$\mathbb{E}(T_{1,n}(p)) = \frac{1}{1 - p_0} = \frac{1}{1 - v_0} = \mathbb{E}(T_{1,n}(v)).$$

For $c \geq 2$, which implies that $n \geq 2$, summing Relation (4) for $k \geq 1$, we get

$$\mathbb{E}(T_{c,n}(p)) - 1 = p_0 \mathbb{E}(T_{c,n}(p)) + \sum_{\ell=1}^n p_\ell \mathbb{E}(T_{c-1,n-1}(p^{(\ell)})).$$

We then obtain

$$(9) \quad \mathbb{E}(T_{c,n}(p)) = \frac{1}{1 - p_0} \left(1 + \sum_{\ell=1}^n p_\ell \mathbb{E}(T_{c-1,n-1}(p^{(\ell)})) \right).$$

We now proceed by recurrence. Suppose that the inequality is true for integer $c - 1$, with $c \geq 2$, i.e., suppose that, for every $n \geq c$, for every $q = (q_1, \dots, q_{n-1}) \in (0, 1)^{n-1}$ with $q_1 + \dots + q_{n-1} \leq 1$, we have

$$\mathbb{E}(T_{c-1,n-1}(q)) \geq \mathbb{E}(T_{c-1,n-1}(v)), \text{ with } v_0 = q_0 = 1 - \sum_{i=1}^{n-1} q_i.$$

Using Relation (8), this implies that

$$\mathbb{E}(T_{c-1,n-1}(p^{(\ell)})) \geq \frac{(n-1)(H_{n-1} - H_{n-c})}{1 - (p_0 + p_\ell)}.$$

From Relation (9), we obtain

$$(10) \quad \mathbb{E}(T_{c,n}(p)) \geq \frac{1}{1 - p_0} \left(1 + (n-1)(H_{n-1} - H_{n-c}) \sum_{\ell=1}^n \frac{p_\ell}{1 - (p_0 + p_\ell)} \right).$$

Observe now that for $\ell = 1, \dots, n$ we have

$$\frac{p_\ell}{1 - (p_0 + p_\ell)} = -1 + \frac{1}{(n-1)r_\ell},$$

where the r_ℓ are given by

$$r_\ell = \frac{1 - (p_0 + p_\ell)}{(n-1)(1 - p_0)}$$

and satisfy $r_1, \dots, r_n \in (0, 1)$ with $r_1 + \dots + r_n = 1$. From Lemma 4, we obtain

$$\sum_{\ell=1}^n \frac{p_\ell}{1 - (p_0 + p_\ell)} = -n + \frac{1}{n-1} \sum_{\ell=1}^n \frac{1}{r_\ell} \geq -n + \frac{n^2}{n-1} = \frac{n}{n-1}.$$

Replacing this value in (10), we obtain, using (8),

$$\mathbb{E}(T_{c,n}(p)) \geq \frac{1}{1 - p_0} (1 + n(H_{n-1} - H_{n-c})) = \frac{n(H_n - H_{n-c})}{1 - p_0} = \mathbb{E}(T_{c,n}(v)),$$

which completes the proof. \square

5. DISTRIBUTION MINIMIZING THE DISTRIBUTION OF $T_{n,n}(p)$

For every $n \geq 1$, $i = 0, 1, \dots, n$ and $k \geq 0$, we denote by $N_i^{(k)}$ the number of coupons of type i collected at instants $1, \dots, k$. It is well-known that the joint distribution of the $N_i^{(k)}$ is a multinomial distribution. More precisely, for every $k \geq 0$ and $k_0, k_1, \dots, k_n \geq 0$ such that $k_0 + k_1 + \dots + k_n = k$, we have

$$(11) \quad \mathbb{P}\{N_0^{(k)} = k_0, N_1^{(k)} = k_1, \dots, N_n^{(k)} = k_n\} = \frac{k!}{k_0!k_1! \dots k_n!} p_0^{k_0} p_1^{k_1} \dots p_n^{k_n}.$$

Recall that the coupons of type 0 do not belong to the collection. For every $\ell = 1, \dots, n$, we easily deduce that, for every $k \geq 0$ and $k_1, \dots, k_\ell \geq 0$ such that $k_1 + \dots + k_\ell \leq k$,

$$\mathbb{P}\{N_1^{(k)} = k_1, \dots, N_\ell^{(k)} = k_\ell\} = \frac{k!}{k_1! \dots k_\ell! (k - (k_1 + \dots + k_\ell))!} p_1^{k_1} \dots p_\ell^{k_\ell} (1 - (p_1 + \dots + p_\ell))^{k - (k_1 + \dots + k_\ell)}.$$

To prove the next theorem, we recall some basic results on convex functions. A function f is said to be convex on an interval I if for every $x, y \in I$ and $\lambda \in [0, 1]$, we have

$$f(\lambda x + (1 - \lambda)y) \leq \lambda f(x) + (1 - \lambda)f(y).$$

Let f be a function defined on an interval I . For every $\alpha \in I$, we introduce the function g_α , defined for every $x \in I \setminus \{\alpha\}$, by

$$g_\alpha(x) = \frac{f(x) - f(\alpha)}{x - \alpha}.$$

It is well-known that f is a convex function on interval I if and only if for every $\alpha \in I$, the function g_α is increasing on $I \setminus \{\alpha\}$. The next result is also known but less popular, so we give its proof.

Lemma 6. *Let f be a convex function on an interval I . For every $x, y, z, t \in I$ with $x < y, z < t$, we have*

$$(t - y)f(z) + (z - x)f(y) \leq (t - y)f(x) + (z - x)f(t).$$

If, moreover, we have $t + x = y + z$, we get

$$f(z) + f(y) \leq f(x) + f(t).$$

Proof. It suffices to apply twice the property that function g_α is increasing on $I \setminus \{\alpha\}$, for every $\alpha \in I$. Since $z < t$, we have $g_x(z) \leq g_x(t)$ and since $x < y$, we have $g_t(x) \leq g_t(y)$. But as $g_x(t) = g_t(x)$ and $g_t(y) = g_y(t)$, we obtain $g_x(z) \leq g_x(t) = g_t(x) \leq g_t(y) = g_y(t)$, which means in particular that

$$\frac{f(z) - f(x)}{z - x} \leq \frac{f(t) - f(y)}{t - y},$$

that is

$$(t - y)f(z) + (z - x)f(y) \leq (t - y)f(x) + (z - x)f(t).$$

The rest of the proof is trivial since $t + x = y + z$ implies that $t - y = z - x > 0$. \square

Theorem 7. *For every $n \geq 1$ and $p = (p_1, \dots, p_n) \in (0, 1)^n$ with $p_1 + \dots + p_n \leq 1$, we have, for every $k \geq 0$,*

$$\mathbb{P}\{T_{n,n}(p') \leq k\} \leq \mathbb{P}\{T_{n,n}(p) \leq k\},$$

where $p' = (p_1, \dots, p_{n-2}, p'_{n-1}, p'_n)$ with $p'_{n-1} = \lambda p_{n-1} + (1 - \lambda)p_n$ and $p'_n = (1 - \lambda)p_{n-1} + \lambda p_n$, for every $\lambda \in [0, 1]$.

Proof. If $\lambda = 1$ then we have $p' = p$ so the result is trivial. If $\lambda = 0$ then we have $p'_{n-1} = p_n$ and $p'_n = p_{n-1}$ and the result is also trivial since the function $\mathbb{P}\{T_{n,n}(p) \leq k\}$ is a symmetric function of p . We thus suppose now that $\lambda \in (0, 1)$.

For every $n \geq 1$ and $k \geq 0$, we have

$$\{T_{n,n}(p) \leq k\} = \{N_1^{(k)} > 0, \dots, N_n^{(k)} > 0\}.$$

We thus get, for $k_1 > 0, \dots, k_{n-2} > 0$ such that $k_1 + \dots + k_{n-2} \leq k$, setting $s = k - (k_1 + \dots + k_{n-2})$,

$$\begin{aligned} \mathbb{P}\{T_{n,n}(p) \leq k, N_1^{(k)} = k_1, \dots, N_{n-2}^{(k)} = k_{n-2}\} \\ = \mathbb{P}\{N_1^{(k)} = k_1, \dots, N_{n-2}^{(k)} = k_{n-2}, N_{n-1}^{(k)} > 0, N_n^{(k)} > 0\} \\ = \sum_{\substack{u \geq 0, v > 0, w > 0, \\ u+v+w=s}} \mathbb{P}\{N_0^{(k)} = u, N_1^{(k)} = k_1, \dots, N_{n-2}^{(k)} = k_{n-2}, N_{n-1}^{(k)} = v, N_n^{(k)} = w\}. \end{aligned}$$

Using Relation (11) and introducing the notation

$$q_0 = \frac{p_0}{p_0 + p_{n-1} + p_n}, \quad q_{n-1} = \frac{p_{n-1}}{p_0 + p_{n-1} + p_n} \text{ and } q_n = \frac{p_n}{p_0 + p_{n-1} + p_n},$$

we obtain

$$\begin{aligned} \mathbb{P}\{T_{n,n}(p) \leq k, N_1^{(k)} = k_1, \dots, N_{n-2}^{(k)} = k_{n-2}\} \\ = \sum_{\substack{u \geq 0, v > 0, w > 0, \\ u+v+w=s}} \frac{k! p_0^u p_1^{k_1} \dots p_{n-2}^{k_{n-2}} p_{n-1}^v p_n^w}{u! k_1! \dots k_{n-2}! v! w!} \\ = \frac{k! p_1^{k_1} \dots p_{n-2}^{k_{n-2}}}{k_1! \dots k_{n-2}!} \sum_{\substack{u \geq 0, v > 0, w > 0, \\ u+v+w=s}} \frac{p_0^u p_{n-1}^v p_n^w}{u! v! w!} \\ = \frac{k! p_1^{k_1} \dots p_{n-2}^{k_{n-2}} (1 - (p_1 + \dots + p_{n-2}))^s}{k_1! \dots k_{n-2}! s!} \sum_{\substack{u \geq 0, v > 0, w > 0, \\ u+v+w=s}} \frac{s!}{u! v! w!} q_0^u q_{n-1}^v q_n^w \\ = \frac{k! p_1^{k_1} \dots p_{n-2}^{k_{n-2}} (1 - (p_1 + \dots + p_{n-2}))^s}{k_1! \dots k_{n-2}! s!} (1 - (q_0 + q_{n-1})^s - (q_0 + q_n)^s + q_0^s). \end{aligned}$$

Note that this relation is not true if at least one of the k_ℓ is zero. Indeed, if $k_\ell = 0$ for some $\ell = 1, \dots, n-2$, we have

$$\mathbb{P}\{T_{n,n}(p) \leq k, N_1^{(k)} = k_1, \dots, N_{n-2}^{(k)} = k_{n-2}\} = 0.$$

Summing over all the k_1, \dots, k_{n-2} such that $k_1 + \dots + k_{n-2} \leq k$, we get

$$(12) \quad \mathbb{P}\{T_{n,n}(p) \leq k\} = \sum_{(k_1, \dots, k_{n-2}) \in E_{n-2}} \frac{k! p_1^{k_1} \dots p_{n-2}^{k_{n-2}} (1 - (p_1 + \dots + p_{n-2}))^s}{k_1! \dots k_{n-2}! s!} (1 - (q_0 + q_{n-1})^s - (q_0 + q_n)^s + q_0^s),$$

where the set E_{n-2} is defined by

$$E_{n-2} = \{(k_1, \dots, k_{n-2}) \in (\mathbb{N}^*)^{n-2} \mid k_1 + \dots + k_{n-2} \leq k\}$$

and \mathbb{N}^* is the set of positive integers.

Note that for $n = 2$, since $p_0 + p_1 + p_2 = 1$, we have $\mathbb{P}\{T_{2,2}(p) \leq k\} = (1 - (p_0 + p_1)^k - (p_0 + p_2)^k + p_0^k)$.

Recall that $p_0 = 1 - (p_1 + \dots + p_n)$. By definition of p'_{n-1} and p'_n , we have, for every $\lambda \in (0, 1)$, $p'_{n-1} + p'_n = p_{n-1} + p_n$. It follows that, by definition of p' ,

$$p'_0 = 1 - (p_1 + \dots + p_{n-2} + p'_{n-1} + p'_n) = 1 - (p_1 + \dots + p_{n-2} + p_{n-1} + p_n) = p_0.$$

Suppose that we have $p_{n-1} < p_n$. This implies, by definition of p'_{n-1} and p'_n , that $p_{n-1} < p'_{n-1}, p'_n < p_n$, that is $q_{n-1} < q'_{n-1}, q'_n < q_n$, where

$$q'_{n-1} = \frac{p'_{n-1}}{p'_0 + p'_{n-1} + p'_n} = \frac{p'_{n-1}}{p_0 + p_{n-1} + p_n} \text{ and } q'_n = \frac{p'_n}{p'_0 + p'_{n-1} + p'_n} = \frac{p'_n}{p_0 + p_{n-1} + p_n}.$$

In the same way, we have

$$q'_0 = \frac{p'_0}{p'_0 + p'_{n-1} + p'_n} = \frac{p_0}{p_0 + p_{n-1} + p_n} = q_0.$$

We thus have $q_0 + q_{n-1} < q'_0 + q'_{n-1}, q'_0 + q'_n < q_0 + q_n$. The function f defined by $f(x) = x^s$ is convex on interval $[0, 1]$ so, from Lemma 6, since $2q_0 + q_{n-1} + q_n = 2q'_0 + q'_{n-1} + q'_n$, we have

$$(13) \quad (q'_0 + q'_{n-1})^s + (q'_0 + q'_n)^s \leq (q_0 + q_{n-1})^s + (q_0 + q_n)^s.$$

Similarly, if $p_n < p_{n-1}$, we have, by definition, $p_n < p'_n, p'_{n-1} < p_{n-1}$, that is $q_n < q'_n, q'_{n-1} < q_{n-1}$ and thus we also have Relation (13) in this case. Using Relation (13) in Relation (12), we get, since $q'_0 = q_0$,

$$\begin{aligned} P\{T_{n,n}(p) \leq k\} &\leq \sum_{(k_1, \dots, k_{n-2}) \in E_{n-2}} \frac{k! p_1^{k_1} \dots p_{n-2}^{k_{n-2}} (1 - (p_1 + \dots + p_{n-2}))^s}{k_1! \dots k_{n-2}! s!} (1 - (q'_0 + q'_{n-1})^s - (q'_0 + q'_n)^s + q_0^s) \\ &= P\{T_{n,n}(p') \leq k\}, \end{aligned}$$

which completes the proof. \square

The function $P\{T_{n,n}(p) \leq k\}$, as a function of p , being symmetric, this theorem can easily be extended to the case where the two entries p_{n-1} and p_n of p , which are different from the entries p'_{n-1} and p'_n of p' , are any $p_i, p_j \in \{p_1, \dots, p_n\}$, with $i \neq j$.

In fact, we have shown in this theorem that for fixed n and k , the function of p , $P\{T_{n,n}(p) \leq k\}$, is a Schur-convex function, that is, a function that preserves the order of majorization. See [5] for more details on this subject.

Theorem 8. *For every $n \geq 1$ and $p = (p_1, \dots, p_n) \in (0, 1)^n$ with $p_1 + \dots + p_n \leq 1$, we have, for every $k \geq 0$,*

$$P\{T_{n,n}(p) > k\} \geq P\{T_{n,n}(v) > k\} \geq P\{T_{n,n}(u) > k\},$$

where $v = (v_1, \dots, v_n)$ with $v_i = (1 - p_0)/n$ and $p_0 = 1 - (p_1 + \dots + p_n)$ and where $u = (1/n, \dots, 1/n)$.

Proof. To prove the first inequality, we apply successively and at most $n - 1$ times Theorem 7 as follows. We first choose two different entries of p , say p_i and p_j such that $p_i < (1 - p_0)/n < p_j$ and next to define p'_i and p'_j by

$$p'_i = \frac{1 - p_0}{n} \text{ and } p'_j = p_i + p_j - \frac{1 - p_0}{n}.$$

With respect to Theorem 7, this leads to write $p'_i = \lambda p_i + (1 - \lambda)p_j$ and $p'_j = (1 - \lambda)p_i + \lambda p_j$, with

$$\lambda = \frac{p_j - \frac{1 - p_0}{n}}{p_j - p_i}.$$

From Theorem 7, the vector p' that we obtain by taking the other entries equal to those of p , i.e., by taking $p'_\ell = p_\ell$, for $\ell = i, j$, is such that

$$P\{T_{n,n}(p) > k\} \geq P\{T_{n,n}(p') > k\}.$$

Note that at this point vector p' has at least one entry equal to $(1 - p_0)/n$, so repeating at most $n - 1$ this procedure, we get vector v .

To prove the second inequality, we use Relation (11). Introducing, for every $n \geq 1$, the set F_n defined by

$$F_n(\ell) = \{(k_1, \dots, k_n) \in (\mathbb{N}^*)^n \mid k_1 + \dots + k_n = \ell\}.$$

For $k < n$, both terms are zero, so we suppose that $k \geq n$. We have

$$\begin{aligned}
P\{T_{n,n}(v) \leq k\} &= P\{N_1^{(k)} > 0, \dots, N_n^{(k)} > 0\} \\
&= \sum_{k_0=0}^{k-n} P\{N_0^{(k)} = k_0, N_1^{(k)} > 0, \dots, N_n^{(k)} > 0\} \\
&= \sum_{k_0=0}^{k-n} \sum_{(k_1, \dots, k_n) \in F_n(k-k_0)} \frac{k!}{k_0! k_1! \dots k_n!} p_0^{k_0} \left(\frac{1-p_0}{n}\right)^{k-k_0} \\
&= \sum_{k_0=0}^{k-n} \binom{k}{k_0} p_0^{k_0} (1-p_0)^{k-k_0} \frac{1}{n^{k-k_0}} \sum_{(k_1, \dots, k_n) \in F_n(k-k_0)} \frac{(k-k_0)!}{k_1! \dots k_n!}.
\end{aligned}$$

Setting $p_0 = 0$, we get

$$P\{T_{n,n}(u) \leq k\} = \frac{1}{n^k} \sum_{(k_1, \dots, k_n) \in F_n(k)} \frac{k!}{k_1! \dots k_n!}.$$

This leads to

$$\begin{aligned}
P\{T_{n,n}(v) \leq k\} &= \sum_{k_0=0}^{k-n} \binom{k}{k_0} p_0^{k_0} (1-p_0)^{k-k_0} P\{T_{n,n}(u) \leq k-k_0\} \\
&\leq P\{T_{n,n}(u) \leq k\} \sum_{k_0=0}^{k-n} \binom{k}{k_0} p_0^{k_0} (1-p_0)^{k-k_0} \\
&\leq P\{T_{n,n}(u) \leq k\},
\end{aligned}$$

which completes the proof. \square

To illustrate the steps used in the proof of this theorem, we take the following example. Suppose that $n = 5$ and $p = (1/16, 1/6, 1/4, 1/8, 7/24)$. This implies that $p_0 = 5/48$ and $(1-p_0)/n = 43/240$. In a first step, taking $i = 4$ and $j = 5$, we get

$$p^{(1)} = (1/16, 1/6, 1/4, 43/240, 19/80).$$

In a second, taking $i = 2$ and $j = 5$, we get

$$p^{(2)} = (1/16, 43/240, 1/4, 43/240, 9/40).$$

In a third step, taking $i = 1$ and $j = 3$, we get

$$p^{(3)} = (43/240, 43/240, 2/15, 43/240, 9/40).$$

For the fourth and last step, taking $i = 5$ and $j = 3$, we get

$$p^{(4)} = (43/240, 43/240, 43/240, 43/240, 43/240) = \frac{43}{48} (1/5, 1/5, 1/5, 1/5, 1/5).$$

6. A NEW CONJECTURE

In this section, we propose a new conjecture stating that the complementary distribution function of $T_{c,n}$ is minimal when the distribution p is equal to the uniform distribution u .

Conjecture. *For every $n \geq 1$, $c = 1, \dots, n$ and $p = (p_1, \dots, p_n) \in (0, 1)^n$ with $p_1 + \dots + p_n \leq 1$, we have, for every $k \geq 0$,*

$$P\{T_{c,n}(p) > k\} \geq P\{T_{c,n}(v) > k\} \geq P\{T_{c,n}(u) > k\},$$

where $v = (v_1, \dots, v_n)$ with $v_i = (1-p_0)/n$ and $p_0 = 1 - (p_1 + \dots + p_n)$ and where $u = (1/n, \dots, 1/n)$.

This new conjecture is motivated by the following facts:

- the result is true for the expectations, see Theorem 5.
- the result is true for $c = n$, see Theorem 8.

- the result is trivially true for $c = 1$ since

$$\mathbb{P}\{T_{1,n}(p) > k\} = \mathbb{P}\{T_{1,n}(v) > k\} = p_0^k \geq 1_{\{k=0\}} = \mathbb{P}\{T_{1,n}(u) > k\}.$$

- the result is true for $c = 2$, see Theorem 9 below.

Theorem 9. *For every $n \geq 2$ and $p = (p_1, \dots, p_n) \in (0, 1)^n$ with $p_1 + \dots + p_n \leq 1$, we have, for every $k \geq 0$,*

$$\mathbb{P}\{T_{2,n}(p) > k\} \geq \mathbb{P}\{T_{2,n}(v) > k\} \geq \mathbb{P}\{T_{2,n}(u) > k\},$$

where $v = (v_1, \dots, v_n)$ with $v_i = (1 - p_0)/n$ and $p_0 = 1 - (p_1 + \dots + p_n)$ and where $u = (1/n, \dots, 1/n)$.

Proof. From Relation (2), we have

$$\mathbb{P}\{T_{2,n}(p) > k\} = -(n-1)p_0^k + \sum_{\ell=1}^n (p_0 + p_\ell)^k$$

and

$$\mathbb{P}\{T_{2,n}(v) > k\} = -(n-1)p_0^k + n \left(p_0 + \frac{1-p_0}{n} \right)^k.$$

For every constant $a \geq 0$, the function $f(x) = (a+x)^k$ is a convex on $[0, \infty[$, so we have, taking $a = p_0$, by the Jensen inequality

$$\left(p_0 + \frac{1-p_0}{n} \right)^k = \left(\frac{1}{n} \sum_{\ell=1}^n (p_0 + p_\ell) \right)^k \leq \frac{1}{n} \sum_{\ell=1}^n (p_0 + p_\ell)^k.$$

This implies that $\mathbb{P}\{T_{2,n}(p) > k\} \geq \mathbb{P}\{T_{2,n}(v) > k\}$.

To prove the second inequality, we define the function $F_{n,k}$ on interval $[0, 1]$ by

$$F_{n,k}(x) = -(n-1)x^k + n \left(x + \frac{1-x}{n} \right)^k.$$

We then have $F_{n,k}(p_0) = \mathbb{P}\{T_{2,n}(v) > k\}$ and $F_{n,k}(0) = \mathbb{P}\{T_{2,n}(u) > k\}$. The derivative of function $F_{n,k}$ is

$$F'_{n,k}(x) = k(n-1) \left[\left(x + \frac{1-x}{n} \right)^{k-1} - x^{k-1} \right] \geq 0.$$

Function $F_{n,k}$ is thus an increasing function, which means that $\mathbb{P}\{T_{2,n}(v) > k\} \geq \mathbb{P}\{T_{2,n}(u) > k\}$. \square

7. APPLICATION TO THE DETECTION OF DISTRIBUTED DENY OF SERVICE ATTACKS

A Deny of Service (DoS) attack tries to progressively take down an Internet resource by flooding this resource with more requests than it is capable to handle. A Distributed Deny of Service (DDoS) attack is a DoS attack triggered by thousands of machines that have been infected by a malicious software, with as immediate consequence the total shut down of targeted web resources (*e.g.*, e-commerce websites). A solution to detect and to mitigate DDoS attacks is to monitor network traffic at routers and to look for highly frequent signatures that might suggest ongoing attacks. A recent strategy followed by the attackers is to hide their massive flow of requests over a multitude of routes, so that locally, these flows do not appear as frequent, while globally they represent a significant portion of the network traffic. The term “iceberg” has been recently introduced to describe such an attack as only a very small part of the iceberg can be observed from each single router. The approach adopted to defend against such new attacks is to rely on multiple routers that locally monitor their network traffic, and upon detection of potential icebergs, inform a monitoring server that aggregates all the monitored information to accurately detect icebergs. Now to prevent the server from being overloaded by all the monitored information, routers continuously keep track of the c (among n) most recent high flows (modelled as items) prior to sending them to the server, and throw away all the items that appear with a small probability p_i , and such that the sum of these small probabilities is modelled by probability p_0 . Parameter c is dimensioned so that the frequency at which all the routers send their c last frequent items is low enough to enable the server to aggregate all of them and to trigger a DDoS

alarm when needed. This amounts to compute the time needed to collect c distinct items among n frequent ones. Moreover, Theorem 5 shows that the expectation of this time is minimal when the distribution of the frequent items is uniform.

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